STAR: SECRET SHARING FOR THRESHOLD AGGREGATION REPORTING

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k = 2



Sometimes known as k-heavy-hitters

THRESHOLD AGGREGATION



Ideal case: No efficient
solutions



N-server aggregation: DPFs, Prio, SMPC



Trusted shuffling: e.g. Prochlo



Approximate: DP, randomised resp.

PREVIOUS WORK



- Emphasis on simplicity and performance
- Well-known cryptography (secret sharing, OPRFs)
- ◇ Orders of magnitude cheaper than state-of-the-art
- ◊ Malicious security
- ◊ Auxiliary data support
- ◇ Open-source rust code: github.com/brave/sta-rs

OVERVIEW OF STAR





Anonymizing proxy (such as Tor, or Oblivious HTTP)

Shamir secret sharing





Symmetric encryption

Oblivious PRF

CRYPTOGRAPHIC TOOLS



THE STAR PROTOCOL

THE PROBLEM	CONSTRUCTION	ANALYSIS	CONCLUSION
Randomn	ess phase		
	[x]	randomness	
	[PRF(sk.x)]	server	

Message phase

$$\diamond (r_1, r_2, r_3) = H(PRF(sk, x))$$

 \diamond **s** = Share(secret=r₁;randomness=r₂), **t** = r₃

$$\diamond$$
 ek = Derive(r₁)

RANDOMNESS & MESSAGE PHASES

THE PROBLEM CONSTRUCTIO	N ANA	LYSIS	CONCLUSIO
Aggregation phase			
[c.s.t]		_	
	aggregation		
~	server		

Steps

- \diamond Group messages based on deterministic tag t
- \diamond If \geq k messages in the group, run share recovery on s and retrieve r_1
- \diamond Derive **ek** from **r**₁
- Output each c to learn (x,aux)

AGGREGATION PHASE

THE	PROBLEM	CONSTRUCTION	ANALYSIS	CONCLUSION
	Malicious	security in	random oracle	model
	[c,s, <mark>t</mark>]		[x]	randomness

Problem: Deterministic tags **Solution:** Randomness server key rotations

aggregation

server



Problem: Randomness DoS
Solution: Clients can verify
randomness (VOPRF)



Problem: Sybil attacks **Solution:** All threshold aggregation schemes vulnerable



Problem: Client identity **Solution:** Proxy messages, e.g. via Tor, or via randomness server using Oblivious HTTP

SECURITY & LEAKAGE

Aggregation runtimes (k \in {0.01%, 0.1%, 1%})



Other costs (per-client)

- Ocmmunication:
 - Aggregation: 233 bytes (+ auxiliary data)
 - Randomness server: 165 bytes
- ◊ VOPRF: < 2ms</p>
- ◊ OHTTP: < 1ms, and approx. 4x communication</p>

PERFORMANCE (256-BIT MEASUREMENTS)

Features

Feature	STAR	Poplar (S&P'21)
Aggregation servers (#)	1	2
Auxiliary data	1	×
Leakage	Tag-based	Prefix-based
Identity-hiding	✓ (OHTTP)	1
Cryptography	Secret-sharing, VOPRFs	Distributed point functions

Headlines (including OHTTP)

- ◊ Computation: 1773x faster
- ◊ Bandwidth: 62.4x smaller
- ◊ Financial: 24x cheaper¹

¹AWS c4.8xlarge, Feb 2022 prices

COMPARISON (STAR & POPLAR)

 Simple, Cheap Privacy-Preserving Threshold Aggregation with k-anonymity

- ◊ Implementations:
 - github.com/brave/sta-rs (Rust)
 - github.com/chris-wood/star-go (Go)
- ◊ IETF standardization: draft-dss-star-02
- ◊ Used at Brave for Web Discovery Project

CONCLUSIONS & ONGOING WORK